Algebraicity of solutions of DDEs with one catalytic variable¹

MathExp (Inria Saclay)

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¹ Joint work with Hadrien Notarantonio. Accepted as a talk at FPSAC23.

Motivating examples

Let $F(t,u) = \sum_{n,k \ge 0} a_{n,k} t^n u^k$ be the generating function of walks in \mathbb{N}^2 which have *n* steps in $\{\nearrow, \searrow\}$ and end at level (height) *k*. One finds:

$$F(t,u) = 1 + tuF(t,u) + t\frac{F(t,u) - F(t,0)}{u}.$$

System case

It follows that:
$$F(t,0) = \frac{1-\sqrt{1-4t}}{2t}$$
. In particular: $F(t,0) \in \overline{\mathbb{Q}(t)}$.

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. In particular: $F(t,0) \in \overline{\mathbb{Q}(t)}$.

■ Modelling [Bonichon, Bousquet-Mélou, Dorbec, Pennarun, 2006] special Eulerian planar orientations gives rise to:

$$\begin{cases} F_1(t,u) = 1 + t \cdot \left(u + 2uF_1(t,u)^2 + 2uF_2(t,1) + u\frac{F_1(t,u) - uF_1(t,1)}{u-1}\right), \\ F_2(t,u) = t \cdot \left(2uF_1(t,u)F_2(t,u) + uF_1(t,u) + uF_2(t,1) + u\frac{F_2(t,u) - uF_2(t,1)}{u-1}\right). \end{cases}$$

Again: $G = F_1(t, 0)$ and $F_2(t, 0)$ are algebraic functions, for example:

$$64t^3G^3 + 2t\left(24t^2 - 36t + 1\right)G^2 - \left(15t^3 - 9t^2 - 19t + 1\right)G + t^3 + 27t^2 - 19t + 1 = 0.$$

Discrete Differential Equations (DDEs) with one catalytic variable

■ The divided difference operator (discrete derivative):

$$\Delta_a: \mathbb{Q}[u]\llbracket t \rrbracket \to \mathbb{Q}[u]\llbracket t \rrbracket,$$

$$F(t,u) \mapsto \frac{F(t,u) - F(t,a)}{u-a}.$$

System case

• Δ_a^j is the *j*-th iteration of Δ_a . Explicitly:

$$\Delta_{a}^{j+1}F(t,u) = \frac{F(t,u) - F(t,a) - (u-a)\partial_{u}F(t,a) - \cdots - \frac{(u-a)^{j}}{j!}\partial_{u}^{j}F(t,a)}{(u-a)^{j+1}}$$

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lacksquare Δ_a^j is the j-th iteration of Δ_a . Explicitly:

$$\Delta_a^{j+1}F(t,u)=\frac{F(t,u)-F(t,a)-(u-a)\partial_uF(t,a)-\cdots-\frac{(u-a)^j}{j!}\partial_u^jF(t,a)}{(u-a)^{j+1}}$$

lacktriangle For polynomials $f(u) \in \mathbb{Q}[u]$ and $Q \in \mathbb{Q}[x,y_1\ldots,y_k,t,u]$ consider the equation

$$F(t,u) = f(u) + t \cdot Q(F(t,u), \Delta_a F(t,u), \dots, \Delta_a^k F(t,u), t, u), \tag{DDE}$$

where $a \in \mathbb{Q}$ (usually 0 or 1) and $k \in \mathbb{N}$ (the order of the **DDE**).

Some history

- In 1960s, Tutte reduced many combinatorial problems to studying **DDEs**. E.g.: [Tutte, 1962] and [Brown, Tutte 1964]
- In 1986 Popescu proved the "General Néron desingularization" [Popescu, 1986].
- [Bandarier, Flajolet 2002]: Universal "Kernel method" for linear **DDEs**.
- In 2006: The unique solution of any **DDE** is an algebraic function and effective method to compute the minimal polynomial [Bousquet-Mélou, Jehanne, 2006].
- In 2015 [Hauser, Rond] organize a conference on "Artin Approximation": Popescu's Theorem also implies algebraicity of **DDEs** (but in a non-effective way).
- [Buchacher, Kauers, 2020]: Linear systems of **DDEs** have algebraic solutions (effective proof).
- [Bostan, Chyzak, Notarantonio, Safey El Din, 2022]: Fast algorithms for order 1.
- New: [Notarantonio, Y., 2022]: Effective proof that systems of **DDEs** have algebraic solutions.

Main result for scalar equations

Theorem (Bousquet-Mélou, Jehanne, 2006)

Let $k \geq 1$, $a \in \mathbb{Q}$ and $Q \in \mathbb{Q}[x, y_1, \dots, y_k, t, u]$, $f(u) \in \mathbb{Q}[u]$. There exists a unique solution $F(t, u) \in \mathbb{Q}[u][\![t]\!]$ of the functional equation

$$F(t,u) = f(u) + t \cdot Q(F(t,u), \Delta_a F(t,u), \dots, \Delta_a^k F(t,u), t, u),$$

and F(t, u) is algebraic over $\mathbb{Q}(t, u)$. Moreover, there exists an algorithm for computing the minimal polynomial of F(t, u).

■ Let E be the polynomial in $x, z_0, \dots, z_{k-1}, t, u$ that is induced by **DDE**:

$$E(u) := E(\underbrace{F(t,u)}_{x},\underbrace{F(t,0)}_{z_0},\ldots,\underbrace{\partial_u^{k-1}F(t,0)}_{z_{k-1}},t,u) = 0.$$

System case

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$$\partial_u F(t, u) \cdot \partial_x E(u) + \partial_u E(u) = 0.$$
 ($\partial_x E(u)$ is the "kernel")

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- Any solution u = U(t) of $\partial_x E(u) = 0$ also implies $\partial_u E(u) = 0$.
- Obtain 3 equations $(E(u) = 0, \partial_x E(u) = 0, \partial_u E(u) = 0)$ for k + 1 + 1 variables.

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- Obtain 3 equations $(E(u) = 0, \partial_x E(u) = 0, \partial_u E(u) = 0)$ for k + 1 + 1 variables.
- If $\partial_x E(u) = 0$ has k distinct solutions U_1, \ldots, U_k , we can consider:

$$\mathcal{S}_{\mathrm{dup}} := egin{cases} E(u_i) = 0, \ \partial_x E(u_i) = 0, \ \partial_u E(u_i) = 0. \end{cases} \quad ext{for } i = 1, \ldots, k. \ egin{cases} rac{F(t, U_i) \leftrightarrow x_i,}{\partial_u^i F(t, 0) \leftrightarrow z_i,} \ U_i \leftrightarrow u_i. \end{cases}$$

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- We find: 3k equations and 3k variables. Hope: the system \mathcal{S}_{dup} is 0-dimensional.
- In that case can use elimination algorithms and find the annihilating polynomial.

Sketch of proof

- Two issues:
 - $\partial_x E(u) = 0$ does not always have k distinct solutions.

System case

2 Is $\mathcal{S}_{\mathrm{dup}}$ really 0-dimensional?

Sketch of proof

- Two issues:
 - 1 $\partial_x E(u) = 0$ does not always have k distinct solutions.
 - 2 Is S_{dup} really 0-dimensional?
- Given the **DDE**

$$F(t,u) = f(u) + t \cdot Q(F(t,u), \Delta_a F(t,u), \dots, \Delta_a^k F(t,u), t, u),$$

consider the perturbed **DDE**_e:

$$G(t, u, \epsilon) = f(u) + t^2 \cdot Q(G(t, u, \epsilon), \Delta_a G(t, u, \epsilon), \dots, \Delta_a^k G(t, u, \epsilon), t^2, u) + t\epsilon^k \Delta^k G(t, u, \epsilon).$$

■ $G(t, u, \epsilon)$ algebraic over $\mathbb{Q}(t, u, \epsilon) \Rightarrow F(t, u)$ algebraic over $\mathbb{Q}(t, u)$.

Lemma 1

Let E_G be the numerator of DDE_{ϵ} . $\partial_x E_G(u) = 0$ has k distinct solutions in $\overline{\mathbb{Q}(\epsilon)}[t^{\frac{1}{k}}]$.

Lemma 2

The ideal $\langle \mathcal{S}_{\mathrm{dup}} \rangle$: $\det(\mathrm{Jac}_{\mathcal{S}_{\mathrm{dup}}})^{\infty}$ is 0-dimensional and contains the minimal poly. of G.

Sketch of proof (Lemma 1)

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■ The equation $\partial_x E_G = 0$ has the form:

$$u^k = \epsilon^k t + t^2 \sum_{i=0}^k u^{k-i} \partial_{y_i} Q(F(t,u),\ldots,\Delta^k F(t,u),t^2,u)$$

Newton's algorithm implies that we find k solutions of the form

$$U_{\ell}(t) = \epsilon \cdot t^{\frac{1}{k}} \cdot \zeta^{\ell} + O(t^{\frac{2}{k}}),$$

for ζ a primitive k-th root of unity.

Sketch of proof (Lemma 2)

Lemma 2

The ideal $\langle \mathcal{S}_{\text{dup}} \rangle$: $\det(\operatorname{Jac}_{\mathcal{S}_{\text{dup}}})^{\infty}$ is 0-dimensional and contains the minimal poly. of G.

■ Recall: $E_G(u) := E_G(F(t, u), F(t, 0), \dots, \partial_u^{k-1} F(t, 0), t, u)$ where E_G describes

$$G(t, u, \epsilon) = f(u) + t^{2} \cdot Q(G(t, u, \epsilon), \Delta_{a}G(t, u, \epsilon), \dots, \Delta_{a}^{k}G(t, u, \epsilon), t^{2}, u) + t\epsilon^{k}\Delta^{k}G(t, u, \epsilon).$$

$$S_{\text{dup}} := \{E_{G}(u_{i}) = 0, \partial_{x}E_{G}(u_{i}) = 0, \partial_{u}E_{G}(u_{i}) = 0 : i = 1, \dots, k\}.$$

Sketch of proof (Lemma 2)

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$$S_{\text{dup}} := \{E_{G}(u_{i}) = 0, \partial_{x}E_{G}(u_{i}) = 0, \partial_{u}E_{G}(u_{i}) = 0 : i = 1, \dots, k\}.$$

■ Application of Hilbert's Nullstellensatz: Lemma 2 holds if the Jacobian matrix $\operatorname{Jac}_{\mathcal{S}_{\operatorname{dup}}} \in \overline{\mathbb{Q}(\epsilon)} \llbracket t^{\frac{1}{\kappa}} \rrbracket^{3k \times 3k}$ of $\mathcal{S}_{\operatorname{dup}}$ is invertible at

$$(x_{1},...,x_{k},u_{1},...,u_{k},z_{0},...,z_{k-1}) = (G(t,U_{1}),...,G(t,U_{k}),U_{1}(t),...,U_{k}(t),G(t,0),...,\partial_{u}^{k-1}G(t,0)) \in \overline{\mathbb{Q}(\epsilon)}[t^{\frac{1}{k}}]^{3k}.$$

Lemma 2

The ideal $\langle S_{\text{dup}} \rangle$: $\det(\operatorname{Jac}_{S_{\text{dup}}})^{\infty}$ is 0-dimensional and contains the minimal poly. of G.

■ Recall: $E_G(u) := E_G(F(t, u), F(t, 0), \dots, \partial_u^{k-1} F(t, 0), t, u)$ where E_G describes

$$G(t, u, \epsilon) = f(u) + t^{2} \cdot Q(G(t, u, \epsilon), \Delta_{a}G(t, u, \epsilon), \dots, \Delta_{a}^{k}G(t, u, \epsilon), t^{2}, u) + t\epsilon^{k}\Delta^{k}G(t, u, \epsilon).$$

$$S_{\text{dup}} := \{E_{G}(u_{i}) = 0, \partial_{x}E_{G}(u_{i}) = 0, \partial_{u}E_{G}(u_{i}) = 0 : i = 1, \dots, k\}.$$

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$$(G(t, U_1), \dots, G(t, U_k), U_1(t), \dots, U_k(t), G(t, 0), \dots, \partial_u^{k-1} G(t, 0)) \in \overline{\mathbb{Q}(\epsilon)} \llbracket t^{\frac{1}{\kappa}} \rrbracket^{3k}.$$

Bousquet-Mélou and Jehanne compute the determinant explicitly:

$$\det(\operatorname{Jac}_{\mathcal{S}_{\operatorname{dup}}}) = r \cdot \prod (\zeta^i - \zeta^j) \cdot \prod^k \left(\partial_x^2 E(U_j) \partial_u^2 E(U_j) - \partial_x \partial_u E(U_j)^2 \right) + O(t^k) \neq 0.$$

Summary of the proof and algorithm in the scalar case

- Strategy of the proof:
 - **I** Given a **DDE** for F(t, u) consider the perturbed **DDE**_{ϵ}:

$$G(t, u, \epsilon) = f(u) + t^2 \cdot Q(G(t, u, \epsilon), \Delta_a G(t, u, \epsilon), \dots, \Delta_a^k G(t, u, \epsilon), t^2, u) + t\epsilon^k \Delta^k G(t, u, \epsilon).$$

and let E_G be the defining polynomial in the the variables $x, z_0, \ldots, z_{k-1}, t, u, \epsilon$.

- 2 Prove that $\partial_x E_G(u) = 0$ has k disctinct solutions u = U(t) in $\overline{\mathbb{Q}(\epsilon)}[t^{\frac{1}{k}}]$.
- 3 Define $S = (E, \partial_x E_G, \partial_u E_G)$ and let S_{dup} be the duplicated system.
- $\textbf{4} \ \, \mathsf{Show that} \,\, \langle \mathcal{S}_{\mathrm{dup}} \rangle : \mathsf{det}(\mathrm{Jac}_{\mathcal{S}_{\mathrm{dup}}})^{\infty} \,\, \mathsf{is} \,\, \mathsf{0}\text{-dimensional by proving:} \,\, \mathrm{Jac}_{\mathcal{S}_{\mathrm{dup}}} \,\, \mathsf{is} \,\, \mathsf{invertible}.$
- Algorithm:
 - **1** Define $E_G \in \mathbb{Q}[x, z_0, \dots, z_{k-1}, t, u, \epsilon]$ as the numerator of of **DDE** $_{\epsilon}$.
 - **2** Compute $\partial_x E$ and $\partial_u E$. Define \mathcal{S}_{dup} in $x_1, \ldots, x_n, u_1, \ldots, u_n, z_0, \ldots, z_{k-1}$.
 - 3 Saturate S_{dup} by adding the equation $m \cdot \det(\operatorname{Jac}_{S_{\text{dup}}}) 1 = 0$ for a variable m.
 - 4 Compute a non-zero element of $\mathcal{S}_{\mathrm{sat}} \cap \mathbb{Q}[z_0,t]$.

New result: extension for systems

Theorem (Notarantonio, Y., 2022)

Let $n, k \geq 1$ be integers and $f_1, \ldots, f_n \in \mathbb{Q}[u], Q_1, \ldots, Q_n \in \mathbb{Q}[y_1, \ldots, y_{n(k+1)}, t, u]$ be polynomials. Set $\nabla^k F := F, \Delta_a F, \ldots, \Delta_a^k F$. Then the system of equations

$$\begin{cases} F_1 = f_1(u) + t \cdot Q_1(\nabla^k F_1, \dots, \nabla^k F_n, t, u), \\ \vdots & \vdots \\ F_n = f_n(u) + t \cdot Q_n(\nabla^k F_1, \dots, \nabla^k F_n, t, u) \end{cases}$$

admits a unique vector of solutions $(F_1, \ldots, F_n) \in \mathbb{Q}[u][[t]]^n$, and all its components are algebraic functions over $\mathbb{Q}(t, u)$.

Theorem (Notarantonio, Y., 2022)

A system of DDEs with one catalytic variable admits a unique vector of solutions $(F_1,\ldots,F_n)\in\mathbb{Q}[u][[t]]^n$, and all its components are algebraic functions over $\mathbb{Q}(t,u)$.

System case

Example (introduced and solved in [Bonichon, Bousquet-Mélou, Dorbec, Pennarun. 2006]):

$$\begin{cases} F_1(t,u) = 1 + t \cdot \big(u + 2uF_1(t,u)^2 + 2uF_2(t,1) + u\frac{F_1(t,u) - uF_1(t,1)}{u-1}\big), \\ F_2(t,u) = t \cdot \big(2uF_1(t,u)F_2(t,u) + uF_1(t,u) + uF_2(t,1) + u\frac{F_2(t,u) - uF_2(t,1)}{u-1}\big). \end{cases}$$

It holds: $G = F_1(t, 0)$ and $F_2(t, 0)$ are algebraic functions. For example:

$$64t^3G^3 + 2t\left(24t^2 - 36t + 1\right)G^2 - \left(15t^3 - 9t^2 - 19t + 1\right)G + t^3 + 27t^2 - 19t + 1 = 0.$$

Generic system case

■ The system

$$\begin{cases} F_1 = f_1(u) + t \cdot Q_1(\nabla^k F_1, \dots, \nabla^k F_n, t, u), \\ \vdots & \vdots \\ F_n = f_n(u) + t \cdot Q_n(\nabla^k F_1, \dots, \nabla^k F_n, t, u) \end{cases}$$

defines *n* polynomial equations $E_1 = 0, \dots, E_n = 0$ in $\mathbb{Q}[x_1, \dots, x_n, z_0, \dots, z_{nk-1}, u, t]$.

System case

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defines n polynomial equations $E_1 = 0, \dots, E_n = 0$ in $\mathbb{Q}[x_1, \dots, x_n, z_0, \dots, z_{nk-1}, u, t]$.

System case

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■ The "derivative" of $(E_1, ..., E_n)$ with respect to u:

$$\begin{pmatrix} \partial_{x_1} E_1 & \dots & \partial_{x_n} E_1 \\ \vdots & \ddots & \vdots \\ \partial_{x_1} E_n & \dots & \partial_{x_n} E_n \end{pmatrix} \cdot \begin{pmatrix} \partial_u F_1 \\ \vdots \\ \partial_u F_n \end{pmatrix} + \begin{pmatrix} \partial_u E_1 \\ \vdots \\ \partial_u E_n \end{pmatrix} = 0.$$
 (1)

■ The system

$$\begin{cases} F_1 = f_1(u) + t \cdot Q_1(\nabla^k F_1, \dots, \nabla^k F_n, t, u), \\ \vdots & \vdots \\ F_n = f_n(u) + t \cdot Q_n(\nabla^k F_1, \dots, \nabla^k F_n, t, u) \end{cases}$$

defines n polynomial equations $E_1 = 0, \ldots, E_n = 0$ in $\mathbb{Q}[x_1, \ldots, x_n, z_0, \ldots, z_{nk-1}, u, t]$.

■ The "derivative" of $(E_1, ..., E_n)$ with respect to u:

$$\begin{pmatrix} \partial_{x_1} E_1 & \dots & \partial_{x_n} E_1 \\ \vdots & \ddots & \vdots \\ \partial_{x_1} E_n & \dots & \partial_{x_n} E_n \end{pmatrix} \cdot \begin{pmatrix} \partial_u F_1 \\ \vdots \\ \partial_u F_n \end{pmatrix} + \begin{pmatrix} \partial_u E_1 \\ \vdots \\ \partial_u E_n \end{pmatrix} = 0.$$

• Observation: (1) and Det = 0 imply P = 0.

Therefore we have: n+2 equations, and n+nk+1 variables.

Generic system case

- Equations: $S = \{E_1(u), \ldots, E_n(u), Det(u), P(u)\}.$
- If Det(u) = 0 has nk distinct solutions in $\overline{\mathbb{Q}}[t^{\frac{1}{k}}]$, then define the system:

$$\mathcal{S}_{\mathsf{dup}} \coloneqq egin{cases} E_1(u_i) = E_2(u_i) = \cdots = E_n(u_i) = 0, \ \mathsf{Det}(u_i) = 0, \ \mathsf{P}(u_i) = 0. \end{cases}$$
 for $i = 1, \ldots, nk$.

System case

- In total $(n+2) \cdot nk = nk(n+2)$ equations.
- Variables: $\underbrace{x_1, \dots, x_{n^2k}}_{F_i(U_i)}, \underbrace{z_0, \dots, z_{nk-1}}_{\partial^j F_i(t,a)}, \underbrace{u_1, \dots, u_{nk}}_{U_i} \Rightarrow n^2k + nk + nk = nk(n+2).$
- \Rightarrow Can hope for a 0-dimensional system.

Outline of the proof for systems in the general case

- **1** Given a system of **DDE**s for $F_i(t, u)$ consider the perturbed system **DDE**_{ϵ}.
- Prove that $\operatorname{Det}(u) = 0$ has nk disctinct solutions u = U(t) in $\overline{\mathbb{Q}(\epsilon)} \llbracket t^{\frac{1}{k}} \rrbracket$.
- Define $S = (E_1, \dots, E_n, Det, P)$ and let S_{dup} be the duplicated system in nk(n+2) variables.
- 4 Show: $\langle S_{\text{dup}} \rangle$: $\det(\operatorname{Jac}_{S_{\text{dup}}})^{\infty}$ is 0-dimensional by proving: $\operatorname{Jac}_{S_{\text{dup}}}$ is invertible.

Sketch of proof: Step 1, deformation

$$\begin{cases} F_1 = f_1(u) + t \cdot Q_1(\nabla^k F_1, \dots, \nabla^k F_n, t, u) \\ \vdots & \vdots \\ F_n = f_n(u) + t \cdot Q_n(\nabla^k F_1, \dots, \nabla^k F_n, t, u) \end{cases}$$

$$\downarrow \downarrow$$

$$\begin{cases} F_{1} = f_{1}(u) + t \cdot Q_{1}(\nabla^{k}F_{1}, \dots, \nabla^{k}F_{n}, t, u), \\ \vdots & \vdots \\ F_{n} = f_{n}(u) + t \cdot Q_{n}(\nabla^{k}F_{1}, \dots, \nabla^{k}F_{n}, t, u) \end{cases}$$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow$$

where $\gamma_{i,i} = i^k$ and $\gamma_{i,i} = t^{\beta}$, and $\alpha \gg \beta \gg 0$ are chosen sufficiently large.

- It holds: $G_i(t, u, \epsilon)$ algebraic over $\mathbb{Q}(t, u, \epsilon) \Rightarrow F_i(t, u)$ algebraic over $\mathbb{Q}(t, u)$.
- **Each** equation for G_i induces a polynomial equation:

$$E_i(\nabla^k G_1, \nabla^k G_2, \ldots, \nabla^k G_n, t, u, \epsilon) = 0, \quad i = 1, \ldots, n.$$

Step 2: Definition of Det(u) and proof of nk distinct solutions

$$\mathsf{Det} := \mathsf{det} \begin{pmatrix} \partial_{\mathsf{x}_1} E_1 & \dots & \partial_{\mathsf{x}_n} E_1 \\ \vdots & \ddots & \vdots \\ \partial_{\mathsf{x}_1} E_n & \dots & \partial_{\mathsf{x}_n} E_n \end{pmatrix}$$

Lemma 1

The equation $\operatorname{Det}(u)=0$ has nk distinct solutions in $\overline{\mathbb{Q}(\epsilon)}\llbracket t^{\frac{1}{\kappa}} \rrbracket$

$$\mathsf{Det} \coloneqq \mathsf{det} \begin{pmatrix} \partial_{\mathsf{x}_1} \mathsf{E}_1 & \dots & \partial_{\mathsf{x}_n} \mathsf{E}_1 \\ \vdots & \ddots & \vdots \\ \partial_{\mathsf{x}_1} \mathsf{E}_n & \dots & \partial_{\mathsf{x}_n} \mathsf{E}_n \end{pmatrix}$$

Lemma 1

The equation $\operatorname{Det}(u)=0$ has nk distinct solutions in $\overline{\mathbb{Q}(\epsilon)}\llbracket t^{\frac{1}{\kappa}} \rrbracket$

■ We have

$$\mathsf{Det}(u) = \mathsf{det} \begin{pmatrix} -u^{m_1} + t\epsilon^k \gamma_{1,1} u^{m_1-k} & \cdots & t\epsilon^k \gamma_{1,n} u^{m_1-k} \\ \vdots & \ddots & \vdots \\ t\epsilon^k \gamma_{n,1} u^{m_n-k} & \cdots & -u^{m_n} + t\epsilon^k \gamma_{n,n} u^{m_n-k} \end{pmatrix} + O(t^{\alpha} u^{M-nk}).$$

- Using $\gamma_{i,i} = i^k$ and $\gamma_{i,j} = t^\beta$, we get $\text{Det} = \prod_{i=1}^n (-u^k + t\epsilon^k j^k) \mod t^{n+1}$.
- Newton's algorithm:

$$U_i = \zeta^{\ell} \cdot j \cdot t^{\frac{1}{k}} + O(t^{\frac{2}{k}}), \quad \text{for } \ell = 1, \dots, k \text{ and } j = 1, \dots, n, \text{ with } \zeta^k = 1.$$

Outline of the proof for systems in the general case

- **1** Given a system of **DDE**s for $F_i(t, u)$ consider the perturbed system **DDE**_{ϵ}.
- 2 Prove that Det(u) = 0 has nk disctinct solutions u = U(t) in $\overline{\mathbb{Q}(\epsilon)} \llbracket t^{\frac{1}{\kappa}} \rrbracket$.
- Define $S = (E_1, \dots, E_n, Det, P)$ and let S_{dup} be the duplicated system in nk(n+2) variables.
- 4 Show: $\langle S_{\text{dup}} \rangle$: $\det(\operatorname{Jac}_{S_{\text{dup}}})^{\infty}$ is 0-dimensional by proving: $\operatorname{Jac}_{S_{\text{dup}}}$ is invertible.

$$\mathsf{Jac}_{\mathcal{S}_\mathsf{dup}} = egin{pmatrix} A_1 & 0 & B_1 \ & \ddots & & dots \ 0 & A_{nk} & B_{nk} \end{pmatrix} \in \overline{\mathbb{Q}(\epsilon)}[[t^{rac{1}{st}}]]^{nk(n+2) imes nk(n+2)},$$

$$A_{i} := \begin{pmatrix} \partial_{x_{1}} E_{1}^{(i)}(U_{i}) & \dots & \partial_{x_{n}} E_{1}^{(i)}(U_{i}) & \partial_{u_{i}} E_{1}^{(i)}(U_{i}) \\ \vdots & \ddots & \vdots & \vdots \\ \partial_{x_{1}} E_{n}^{(i)}(U_{i}) & \dots & \partial_{x_{n}} E_{n}^{(i)}(U_{i}) & \partial_{u_{i}} E_{n}^{(i)}(U_{i}) \\ \partial_{x_{1}} \operatorname{Det}^{(i)}(U_{i}) & \dots & \partial_{x_{n}} \operatorname{Det}^{(i)}(U_{i}) & \partial_{u_{i}} \operatorname{Det}^{(i)}(U_{i}) \\ \partial_{x_{1}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) & \partial_{u_{i}} \operatorname{Det}^{(i)}(U_{i}) \\ \partial_{x_{1}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) & \partial_{u_{i}} P^{(i)}(U_{i}) \\ \partial_{x_{1}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) \\ \partial_{x_{1}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) \\ \partial_{x_{1}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) \\ \partial_{x_{1}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) \\ \partial_{x_{1}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) \\ \partial_{x_{1}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) \\ \partial_{x_{1}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) \\ \partial_{x_{1}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) \\ \partial_{x_{1}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) \\ \partial_{x_{1}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) \\ \partial_{x_{1}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) \\ \partial_{x_{1}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) \\ \partial_{x_{1}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) \\ \partial_{x_{1}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) \\ \partial_{x_{1}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) \\ \partial_{x_{1}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) \\ \partial_{x_{1}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) & \dots & \partial_{x_{n}} P^{(i)}(U_{i}) \\ \partial_{x_{1}} P^{(i)}(U_{i}) & \dots & \partial_{x$$

$$\mathsf{Jac}_i(u) \in \mathbb{O}(\epsilon)[u][[t]]^{(n+1)\times (n+1)} \text{ and } \Lambda \in \overline{\mathbb{O}(\epsilon)}[[t^{\frac{1}{\epsilon}}]]^{nk\times nk}$$

Method: Analyze the (lowest) valuation in t to show non-vanishing.

■ [Bonichon, Bousquet-Mélou, Dorbec, Pennarun, 2006] consider and solve:

$$\begin{cases} F_1(t,u) = 1 + t \cdot \left(u + 2uF_1(t,u)^2 + 2uF_2(t,1) + u\frac{F_1(t,u) - uF_1(t,1)}{u-1}\right), \\ F_2(t,u) = t \cdot \left(2uF_1(t,u)F_2(t,u) + uF_1(t,u) + uF_2(t,1) + u\frac{F_2(t,u) - uF_2(t,1)}{u-1}\right). \end{cases}$$

System case

Special Eulerian planar orientations

We get polynomial equations

$$\begin{cases} E_1 = (1-x_1) \cdot (u-1) + t \cdot (2u^2x_1^2 - u^2z_0 + 2u^2z_1 - 2ux_1^2 + u^2 + ux_1 - 2uz_1 - u), \\ E_2 = x_2 \cdot (1-u) + t \cdot (2u^2x_1x_2 + u^2x_1 - 2ux_1x_2 - ux_1 + ux_2 - uz_1). \end{cases}$$

System case

Then define

$$\begin{cases} \text{Det} = (4tu^2x_1 - 4tux_1 + tu - u + 1)(2tu^2x_1 - 2tux_1 + tu - u + 1), \\ P = -2tx_1x_2 - tx_1 + tx_2 - tz_1 - x_2 + P_1 \cdot u + P_2 \cdot u^2 + P_3 \cdot u^3, \end{cases}$$

- $\mathcal{S}_{\text{dup}} =$ $(E_1(x_1, x_2, z_0, z_1, u_1), E_2(x_1, x_2, z_0, z_1, u_1), Det(x_1, x_2, z_0, z_1, u_1), P(x_1, x_2, z_0, z_1, u_1), (E_1(x_3, x_4, z_0, z_1, u_2), E_2(x_3, x_4, z_0, z_1, u_2), Det(x_3, x_4, z_0, z_1, u_2), P(x_3, x_4, z_0, z_1, u_2))$
- Compute a generator of $\langle S_{\text{dup}}, m \cdot (u_1 u_2) 1 \rangle \cap \mathbb{Q}[z_0, t]$.

Special Eulerian planar orientations in Maple

```
E1 := numer(-x1+1 + 2*t*u*x1^2 + 2*t*u*z1 + t*u*(-u*z0+x1)/(u-1)+t*u);
E2 := numer(-x^2 + 2*t*u*x^1*x^2 + t*u*x^1 + t*u*z^1 + t*u*(-u*z^1+x^2)/(u-1)):
Jac := Matrix([[diff(E1, x1), diff(E1, x2)], [diff(E2, x1), diff(E2, x2)]]);
Det := LinearAlgebra[Determinant](Jac);
Pm := Matrix([[diff(E1, x1), diff(E2, x1)], [diff(E1, u), diff(E2, u)]]);
P := LinearAlgebra[Determinant](Pm);
S := [E1, E2, det, P];
S1 := op(subs(x1=x1, x2=x2, u=u1, S);
S2 := op(subs(x1=x3,x2=x4,u=u2,S));
Sdup := [S1.S2. m*(u1 - u2) - 1]:
G := polynomial_elimination(Sdup, z0, t);
  (z0-1)(2tz0+t-1)(64t^3z0^3+48t^3z0^2-15t^3z0-72t^2z0^2+t^3+9t^2z0+\cdots)\cdots
                                                                            21/22
```

Summary and conclusion

■ Systems of DDEs with one catalytic variable have algebraic solutions.

■ There exists an algorithm for finding minimal polynomials of such solutions.

 Currently ongoing work on improving the efficiency and effective handling of more catalytic variables.